

New Slide Attacks on Almost Self-Similar Ciphers

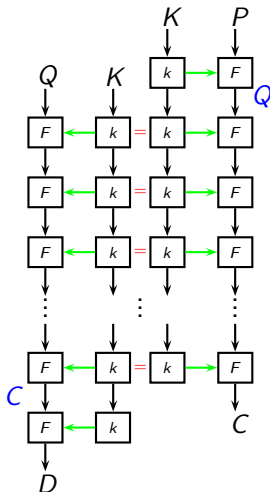
Orr Dunkelman, Nathan Keller,
Noam Lasry, Adi Shamir

May 21th, 2019



Slide Attacks [BW99]

- ▶ Presented by Biryukov and Wagner in 1999
- ▶ Can be applied to ciphers with the same keyed permutation
- ▶ Independent of the number of rounds of the cipher



Slide Attacks [BW99] (cont.)

- ▶ Slid pair satisfies

$$\begin{cases} Q = f_k(P), \\ D = f_k(C), \end{cases} \quad (1)$$

- ▶ Slide attacks:
 - ▶ Find such a slid pair,
 - ▶ Use slid pair to extract key.



Extensions and Generalizations

- ▶ Slide with twist [BW00]
- ▶ Advanced slide [BW00]
- ▶ Chains [F01]
- ▶ Slidex [DKS12]
- ▶ Reflection [K08]
- ▶ Quantum [B+18]



Applications

- ▶ 1K-DES, 2K-DES, 4K-DES ([BW99,BW00])
- ▶ 3K-DES ([B+17])
- ▶ 1K-AES ([B+17])
- ▶ KeeLoq ([I+08,C+08])
- ▶ FF3 ([DV17,HMT19])
- ▶ ...



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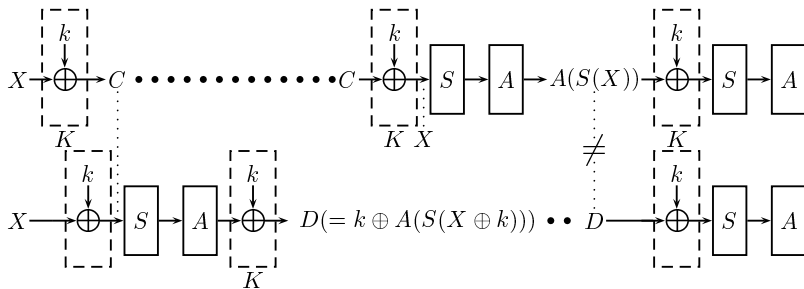
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Problem: in AES the last round is different!



Last Round Function \Rightarrow No Slid Chains



Overcoming the Last Round

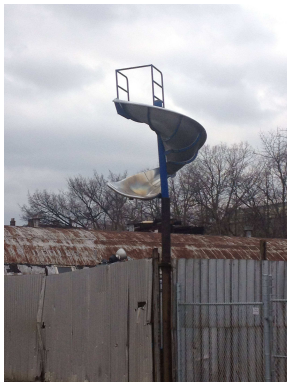
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Introducing **4** new slide techniques:

- ▶ Slid Sets
- ▶ Hypercube of slid pairs
- ▶ Suggestive plaintext structures
- ▶ Substitution slide

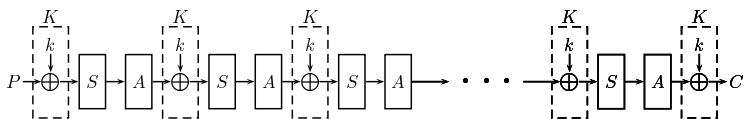


Slid Sets

- ▶ Two sets of λ -structures $\{\mathcal{P}\}$ and $\{\mathcal{Q}\}$ such that

$$f_k(\{\mathcal{P}\}) = \{\mathcal{Q}\}$$

- ▶ Cool detection techniques!
- ▶ Can be used to attack 2K-AES with complexity 2^{68}
- ▶ General 2-KSA — $2^{(n+s)/2}$
- ▶ Can be used to attack 1K-AES with secret S-boxes with complexity $2^{70.3}$



Results

Cipher	Technique	Complexity (general)		AES-like	
		Data/Memory	Time	Data/Memory	Time
Known S-Boxes					
1-KSAf	Slide [B+17]	$2^{n/2}$ (KP)	$2^{n/2}$	2^{64} (KP)	2^{64}
1-KSA _t	Suggestive str.	$3 \cdot 2^{n/2}$ (CP)	$4 \cdot 2^{n/2}$	$2^{65.6}$ (CP)	2^{66}
1-KSA _t	Sub. slide	$2^{n/2}$ (KP)	$2^{3n/4}$	2^{64} (KP)	2^{96}
2-KSAf	Slid sets	$2^{(n+s)/2+1}$ (CP)	$2^{(n+s)/2+1}$	2^{69} (CP)	2^{69}
2-KSAf	Slide + Key Guessing	$(n/s)2^{n/2}$ (CP)	$2^{n/2+s}$	2^{68} (CP)	2^{72}
2-KSA _t †	Slid sets	$2^{(n+m)/2+1}$ (CP)	$\max\{2^{(n+m)/2+1}, 2^{2m}\}$	2^{78} (CP)	2^{78}
3-KSA _f †	Slid sets	$2^{(n+m)/2+1}$ (CP)	$\max\{2^{(n+m)/2+1}, 2^{2m}\}$	2^{81} (CP)	2^{81}
Secret S-Boxes					
1-KSAf	Slid sets	$1.17\sqrt{s}2^{(n+s)/2}$ (CP)	$1.17\sqrt{s}2^{(n+s)/2}$	$2^{70.3}$ (CP)	$2^{70.3}$
1-KSAf	Hypercube	$\sqrt{s}2^{n/2+s(s+3)/4+1}$ (CP)	$\sqrt{s}2^{n/2+s(s+3)/4+1}$	2^{88} (CP)	2^{88}

KP – Known Plaintext; CP – Chosen Plaintext; For AES-like $n = 128, s = 8$

† – this version has incomplete diffusion layer, m denotes the “word” size of the linear operation.

‡ – the memory complexity of this attack is 2^{47} .

Thank you for your Attention!

<https://eprint.iacr.org/2019/059>



A Formal Complaint: Wrongful Rejection from Rump Session

Orr Dunkelman

May 21th, 2019



Facts

- ▶ Eran Lambooj and me wanted to give another rump session presentation

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- ▶ Eran Lambooj and me wanted to give another rump session presentation
- ▶ As you can see, it is not in the schedule
- ▶ The fact that **we missed** the deadline by **13 hours and 27 minutes** seems to be of little relevance!

(Rump) Session Hijacking Attack

- ▶ New attack against timing constraints!

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- ▶ CaML + (rump session talk) MR + Time-Travel attack

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- ▶ New attack against timing constraints!
- ▶ CaML + (rump session talk) MR + Time-Travel attack
- ▶ Also works in the “wrong time zone” model



Thank you for your Support!



A Practical Cryptanalysis Competition

A Cycling Approach

Orr Dunkelman and Eran Lambooj

Eurocrypt, Rumpsession, 2019



What do Cryptographers like the Most?

What do Cryptographers like the Most?

Blockchain

What do Cryptographers like the Most?

BlockKette

Free Drinks

What do Cryptographers like the Most?

BlockKette

Competitions!

~~Freie Getränke~~

We present

The practical LWC competition

Yellow



Most broken
ciphers.

Green



Most points.

Polka dot



Highest broken
rounds.

What can **you** do?

- ▶ Go to the competition website:
`https://cryptanex.hideinplainsight.io/lwc/`
- ▶ Compute key
- ▶ Celebrate
- ▶ Submitters - you can help us (contact Eran)